

QR-factorizations:

Two vectors $x, y \in \mathbb{R}^n$ are orthogonal if $x^T y = \sum_{i=1}^n x_i y_i = 0$

A matrix $Q \in \mathbb{R}^{n \times n}$, is orthogonal if $Q^T Q = I_n$

So Q is invertible with $Q^{-1} = Q^T$.

Examples:

1) The identity matrix: $Q = I$

2) Permutations: $Q = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$

3) Rotations: $Q = \begin{bmatrix} \cos(\theta) & -\sin(\theta) \\ \sin(\theta) & \cos(\theta) \end{bmatrix}$

4) Reflections: $v \in \mathbb{R}^n$

$$Q = I - 2 \frac{v \cdot v^T}{v^T v}$$

Properties:

• ($m = n$) If Q_1, Q_2 are orthogonal, so is $Q = Q_1 Q_2$.

Proof: $Q^T Q = (Q_1 Q_2)^T (Q_1 Q_2) = Q_2^T \underbrace{Q_1^T Q_1}_{I_n} Q_2 = Q_2^T Q_2 = I_n$

• $\|Q \cdot x\|_2 = \|x\|_2$.

Proof: $\|Qx\|_2^2 = x^T Q^T Q x = x^T x = \|x\|_2^2$.

NB! This holds even if $Q \in \mathbb{R}^{n \times m}$, $Q^T Q = I_m$.

Corollary: ($m = n$) $|\lambda(Q)| = 1$.

Some applications:

• If $A = A^T$ then $Q^T A Q = \text{diag}\{\lambda_1, \dots, \lambda_m\}$

• QR-factorization: $A = Q \cdot R$, R upper triangular.

• Singular value decomposition: $A = U \Sigma V$, U, V orthogonal, Σ diagonal.

QR-factorizations

Applications

• Least square problems:

Given the overdetermined system

$$Ax = b, \quad A \in \mathbb{R}^{n \times m}, \quad n > m.$$

Find x s.t. $\|b - Ax\|_2$ is as small as possible.
This is the solution of the normal equation

$$A^T A x = A^T b$$

With $A = \tilde{Q} \cdot \tilde{R} \Rightarrow \tilde{R}^T \tilde{R} x = \tilde{R}^T \tilde{Q} b \Rightarrow \tilde{R} x = \tilde{Q}^T b$

• QR-iterations for finding eigenvalues:

Given $A \in \mathbb{R}^{n \times n}$

for $k = 0, 1, 2, \dots$
Find Q, R s.t. $A^{(k)} = QR$
 $A^{(k+1)} = R \cdot Q$

end

If $|\lambda_1| > |\lambda_2| > \dots > |\lambda_n|$ then

$$A^{(k)} \xrightarrow{k \rightarrow \infty} \begin{bmatrix} \lambda_1 & x & x & \dots & x \\ & \lambda_2 & & & \vdots \\ & & \dots & & x \\ 0 & & & & \lambda_n \end{bmatrix},$$

Algorithm: How to compute Q and R.

Gram-Schmidt:

Let $A = [a_1, \dots, a_m]$, a_i is column i of A

Assume A has full rank, a_i are linear independent.

• Let $q_1 = a_1 / \|a_1\|_2$

• Find $v_2 = a_2 - r_{12} q_1$ so that $v_2 \perp q_1$

that $v_2^T q_1 = a_2^T q_1 - r_{12} \underbrace{q_1^T q_1}_{=1} = 0 \Rightarrow r_{12} = a_2^T q_1$

and let $q_2 = \frac{v_2}{\|v_2\|_2}$.

Assume we have found q_1, \dots, q_{j-1} , based on a_1, a_2, \dots, a_{j-1} . Then, find v_j

$$v_j = a_j - \sum_{i=1}^{j-1} r_{ij} q_i \quad \text{s.t. } v_j \perp q_l, \quad l=1, 2, \dots, j-1$$

$$\Rightarrow v_j^T q_l = a_j^T q_l - \sum_{i=1}^{j-1} r_{ij} \underbrace{q_i^T q_l}_{=\delta_{il}} = a_j^T q_l - r_{lj} = 0$$

$$\Rightarrow r_{lj} = a_j^T q_l$$

Which together give us

Gram-Schmidt algorithm

Given $A \in \mathbb{R}^{n \times m}$ $n \geq m$ with full rank.

for $j=1, 2, \dots, m$

$$v_j = a_j - \sum_{i=1}^{j-1} r_{ij} q_i, \quad r_{ij} = a_j^T q_i$$

$$q_j = v_j / \|v_j\|_2, \quad r_{jj} = \|v_j\|_2$$

end

As long as a_1, a_2, \dots, a_m are linear independent, $v_j \neq 0$ and the algorithm will continue.
If $v_j = 0$, then $a_j \in \text{span}\{a_1, \dots, a_{j-1}\}$

Notice that

$$v_j = r_{jj} \cdot q_j = a_j - \sum_{i=1}^{j-1} r_{ij} q_i$$

$$\Rightarrow a_j = \sum_{i=1}^j r_{ij} q_i$$

so, with

$$\tilde{Q} = [q_1, \dots, q_m] \quad \text{and} \quad R = \begin{bmatrix} r_{11} & \dots & r_{1m} \\ 0 & \ddots & \vdots \\ & & r_{mm} \end{bmatrix} \in \mathbb{R}^{m \times m}$$

we have $\tilde{A} = \tilde{Q} \cdot \tilde{R}$, upper triangular.

This is called the reduced QR-factorization.

Each $A \in \mathbb{R}^{n \times m}$ of full column rank has a unique, reduced QR-factorization

$$A = \tilde{Q} \tilde{R} \quad \text{with } r_{ii} > 0$$

We can extend \tilde{Q} : $Q = [\hat{Q}, q_{m+1}, \dots, q_n] \in \mathbb{R}^{n \times n}$

Where q_{m+1}, \dots, q_n are arbitrary chosen orthogonal vectors, and

$$R = \begin{bmatrix} \tilde{R} \\ 0 \end{bmatrix} \in \mathbb{R}^{n \times m}$$

If A do not have full column rank, then some $r_j = 0 \Rightarrow r_{jj} = 0$. Then pick some q_j orthogonal to the previous ones.
In fact

Every $A \in \mathbb{R}^{n \times m}$, $n \geq m$ has a QR-factorization

$$A = Q \cdot R$$

where Q is orthogonal and $R \in \mathbb{R}^{n \times m}$ is upper triangular.

But: Gram-Schmidt orthogonalization is numerical unstable.

Householder triangulation.

This is an alternative way to construct the QR-factorization. The main tool is the

Householder reflection:

Given some $v \in \mathbb{R}^n$.
Then

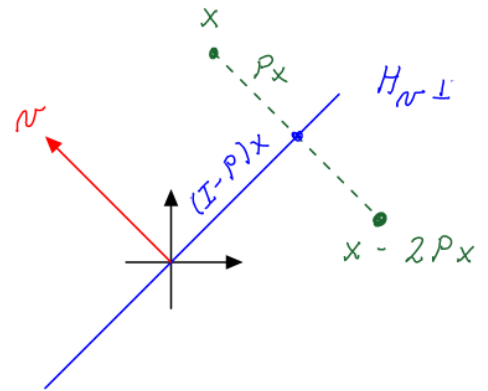
$$P = \frac{vv^T}{\|v\|_2^2}$$

is an orthogonal projection onto $H = \text{span}\{v\}$.

So

$$Q = I - 2P = I - 2 \frac{vv^T}{\|v\|_2^2}$$

is an orthogonal reflection over the hyperplane $\{z \in \mathbb{R}^n, z^T v = 0\} = H_{v^\perp}$.
Moreover $Q = Q^T = Q^{-1}$.



In the following, assume $\|v\|_2 = 1$, so our Householder reflection is

$$Q = I - 2vv^T.$$

Also, notice that

$$Q \cdot x = x - 2 \overbrace{vv^T}^{\mathbb{R}} \cdot x = x - 2v^T x \cdot v.$$

So we never store Q , only v .

Application

Given $x \in \mathbb{R}^n$. Find Q (or v) s.t.

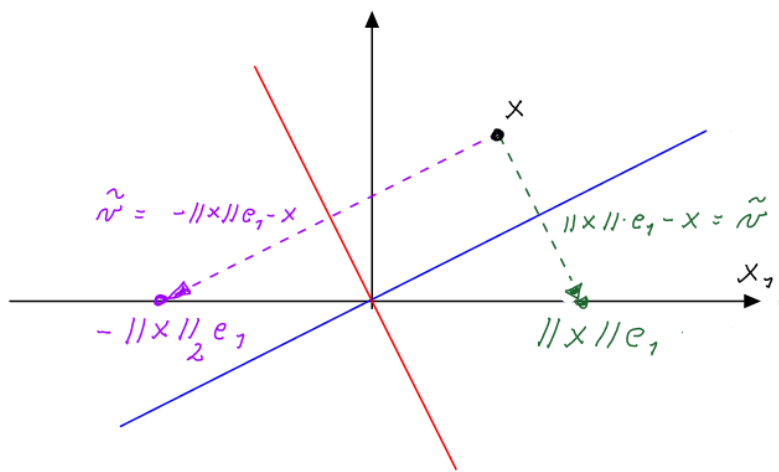
$$Q \cdot x = \pm \begin{bmatrix} \|x\|_2 \\ 0 \\ \vdots \\ 0 \end{bmatrix} = \pm \|x\|_2 \cdot e_1, \quad e_1 = [1, 0, \dots, 0]^T$$

$$\Rightarrow (I - 2 \frac{vv^T}{v^T v}) x = \pm \|x\|_2 \cdot e_1$$

$$\Rightarrow v = \frac{v^T v}{v^T x} (x \mp \|x\|_2 \cdot e_1)$$

NB! We only want the direction of v .

Choose v such that $\|v\|_2 = 1$. Moreover, if we are free to choose the sign, choose the sign of x_1 , to avoid subtraction of almost equal numbers.



This figure is supposed to illustrate the situation.

In conclusion: Choose

$$\tilde{v} = x + \text{sign}(x_1) \cdot \|x\| \cdot e_1, \quad v = \frac{\tilde{v}}{\|\tilde{v}\|_2}$$

Given $A \in \mathbb{R}^{n \times m}$, with full column rank.
We want to find Q_1, Q_2, \dots so that

$$Q_1 \cdot A = \begin{bmatrix} x & x & \dots & x \\ 0 & x & & x \\ 0 & x & & x \\ \vdots & & & \vdots \\ 0 & x & & x \end{bmatrix}, \quad Q_2 Q_1 A = \begin{bmatrix} x & x & x & \dots & x \\ 0 & x & x & \dots & x \\ 0 & 0 & x & \dots & x \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & x & \dots & x \end{bmatrix}$$

Until

$$Q_m Q_{m-1} \dots Q_1 A = R \quad (\text{upper triangular})$$

$$\text{Then } A = Q_1 Q_2 \dots Q_m \cdot R = QR.$$

Householder QR-factorization

$$R = A$$

for $k = 1:m$

$$x = R_{k:n,k}$$

$$v_k = \text{sign}(x_1) \cdot \|x\|_2 \cdot e_1 + x$$

$$v_k = v_k / \|v_k\|_2$$

$$R_{k:n,k:m} = R_{k:n,k:m} - 2v_k (v_k^T R_{k:n,k:m})$$

end.

Notice that $v_k \in \mathbb{R}^{n-k+1}$, and Q_k , as used above is

$$Q_k = \begin{pmatrix} I_{k-1} & 0 \\ 0 & I_{n-k+1} - 2v_k v_k^T \end{pmatrix}$$

The algorithm produce R but not Q .
Usually, we do not need Q itself, only the operation it perform.

For example:

$$Ax = b \quad A = QR \quad \text{nonsingular}$$

$$\Rightarrow QRx = b \Rightarrow Rx = Q^T b = Q_1, Q_2, \dots, Q_n b$$

The computation of $y = Q^T b$ is then

$$y = b$$

for $k = 1:n$

$$y_{k:n} = y_{k:n} - 2N_k(N_k^T y_{k:n})$$

end