

# TMA4155 — Solutions to exercise 4

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## Task 1 — MACs

**a.** We know that  $f(k, x) = t$ . This means that we know one thing about the function  $\tilde{f}(k, \cdot)$ , namely that  $\tilde{f}(k, x) = t_0$ .

Note that for a message  $x||y$  consisting of two blocks  $x$  and  $y$ , the CBC-MAC  $f(k, x||y) = \tilde{f}(k, \tilde{f}(k, x) \oplus y)$ . If we can force the input to the outer evaluation to be  $x$ , then we know that the resulting output will be  $t_0$ . So if we set  $y = x \oplus t_0$ , we get that

$$f(k, x||y) = \tilde{f}(k, \tilde{f}(k, x) \oplus y) = \tilde{f}(k, t_0 \oplus (x \oplus t_0)) = \tilde{f}(k, x) = t_0.$$

**b.** Let  $m' \in \{0, 1\}^n$  be arbitrary. Then

$$t' = f(k, m||m') = h(k||m||m') = h'(h'(h'(0, k), m), m').$$

If we know  $t = f(k, m) = h'(h'(0, k), m)$ , we can compute

$$t' = h'(t, m').$$

## Task 2 — RSA cryptosystem

**a.** We are given  $p = 5$ ,  $q = 7$  and  $e = 5$ . Then  $n = pq = 35$ , and  $d$  should be an inverse to  $e = 5$  modulo  $(p - 1)(q - 1) = 4 \cdot 6 = 24$ . It is easy to see that  $5 \cdot 5 \equiv 25 \equiv 1 \pmod{24}$ , so  $d = 5$  is ok.

This gives the public key  $pk = (35, 5)$  and the secret key (private key)  $sk = (35, 5)$ .

**b.**  $6^5 \equiv 6^{2^2} \cdot 6 \equiv 36^2 \cdot 6 \equiv 1^2 \cdot 6 \equiv 6 \pmod{35}$ .

**c.**  $17^5 \equiv 17^{2^2} \cdot 17 \equiv 9^2 \cdot 17 \equiv 11 \cdot 17 \equiv 12 \pmod{35}$ .

### **Task 3 — Greatest common divisor**

a.  $\gcd(153, 157) = 1$ .

b.  $\gcd(153, 187) = 17$ .

c.  $\gcd(10005, 1101) = 3$ .

d.  $\gcd(10005, 1103) = 1$ .